

Data

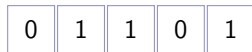
AP Computer Science Principles ■ Unit 2

AP Computer Science Principles

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What we will cover

- 1 Binary numbers
- 2 Data compression
- 3 Extracting information
- 4 Using programs with data



bits

1

Binary

Everything is bits

A **bit** 位 is a single 0 or 1. **Binary** 二进制 (base 2) uses place values that are **powers of two** 2 的幂.

128	64	32	16	8	4	2	1	
0	0	1	0	1	1	0	1	= 45

$32 + 8 + 4 + 1 = 45$. Read binary $\rightarrow$ decimal by adding the place values where a bit is 1.

place value	128	64	32	16	8	4	2	1
bits	1	0	0	1	0	1	1	0

$$150 = 128 + 16 + 4 + 2 \Rightarrow 10010110$$

Fixed bits, overflow & sampling

Convert & overflow

1011_2 : place values 8, 4, 2, 1; bits at 8, 2, 1 give $8 + 2 + 1 = 11$.

With n bits you can make 2^n values, 0 to $2^n - 1$. Exceeding it is an **overflow error** 溢出错误 (4 bits max 15; $15 + 1$ overflows).

Analog 模拟 signals become **digital** 数字 by **sampling** 采样 – many quick measurements. Eight bits group into one **byte** 字节.

A finite number of bits has two consequences

the cause

a computer has a **finite** number of bits, so it can represent only a limited range of values

overflow error 溢出错误

a number **too large** for the available bits cannot be stored correctly

round-off error 舍入错误

numbers with **decimals** can only be **approximated**, because infinitely many real values must map onto finitely many bit patterns

why it matters

$0.1 + 0.2$ does not print as 0.3 in most languages – and that is not a bug in the language

Both are consequences of the same fact. Say “finite bits” first and the explanation writes itself.

Where a number loses precision

Overflow

A value too large for the bits allotted wraps around or errors – the count that goes negative.

Round-off error

Round-off (rounding) error 舍入误差: a real number stored in finite bits is stored *approximately*, so $0.1 + 0.2 \neq 0.3$ exactly.

Accumulation

Round the same way a million times and the error compounds – which is why money is stored in integer cents, not floats.

The shared cause

Both are the same limit: a fixed number of bits cannot represent unlimited values.

An answer that names the limit – “only 2^{32} distinct values exist” – explains both errors at once.

2

Compression

Data compression

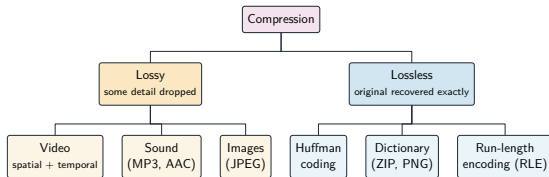
Compression 压缩 stores the same data in **fewer bits** – less space, faster transmission.

Lossless 无损

every bit restored exactly (text, code)

Lossy 有损

discards unnoticed data;
much smaller (photos, music)



Lossy trades **size 大小** against **quality 质量**. Many methods remove **redundancy 冗余** – e.g. “100 white pixels” → “white, 100 times”.

Two kinds of compression, and when each is right

compression 压缩

reducing the number of bits needed to store or send data

lossless compression 无损压缩

the original is **perfectly recoverable** – PNG, ZIP, FLAC. Use it for text, code and anything where a changed bit matters

lossy compression 有损压缩

discards data judged unnoticeable, and it **cannot be recovered** – JPEG, MP3, MP4

the trade

lossy achieves far smaller files; lossless guarantees fidelity

the rule

never use lossy on a program's source, a spreadsheet, or a medical scan

The exam asks which is appropriate, not which is better. Match the choice to **whether losing detail is acceptable**.

3

Extracting information

Extracting information from data

Programs turn raw **data sets** 数据集 into **information** 信息, finding **patterns** 模式 and **trends** 趋势.

Filtering 筛选

keep the rows you want

Cleaning 清洗

fix duplicates, blanks, errors

Beware: a **correlation** 相关 never proves **causation** 因果, and **bias** 偏差 in how data was collected can mislead. **Meta-data** 元数据 & **visualization** 可视化 help.



*Hard-disk platters and head: data compressed
and stored as magnetic patterns*

From data to information

data 数据 into **information** 信息

data becomes useful when patterns, trends and answers are extracted from it

filtering 过滤

keeping only the rows that meet a condition

cleaning 清洗

removing errors, duplicates and incomplete records before analysing

visualizing 可视化

a chart makes a pattern visible that a table hides

the caution

large data sets reveal **correlation**, which is not causation – and they raise **privacy** 隐私 questions about the people in them

Programs process data at scales no human can. That is the benefit *and* the reason the privacy question is on this paper.

4

Using programs

Using programs with data

A program **processes** 处理 data: records in, new information out.

Iterate to summarise

A loop uses **iteration** 迭代 to visit each **element** 元素: summing 70, 85, 90, 60, 95 gives 400; $\div 5 = 80$ (the average).

A **filter** 过滤器 keeps a **subset** 子集 that meets a condition; programs also **combine** 合并 sources, **transform** 转换 values, and **clean** 清洗 first – dirty data gives a wrong answer.

Worked example – choosing a compression

**Which compression for a source-code file?
For a holiday photo?**

Code → **lossless**. Every bit matters – a single changed character breaks the program. Restoring the **exact** original is non-negotiable. **Photo** → **lossy**. A small quality loss is invisible, and the file gets **much** smaller.

The trade is **size** against **fidelity**. Lossy loss is **permanent** – you can never get the discarded data back.

Key points

Binary place values are powers of two; watch **overflow**

Lossless keeps every bit; **lossy** trades quality for size

Correlation is not **causation**; clean data first

Programs **iterate**, **filter** & transform data

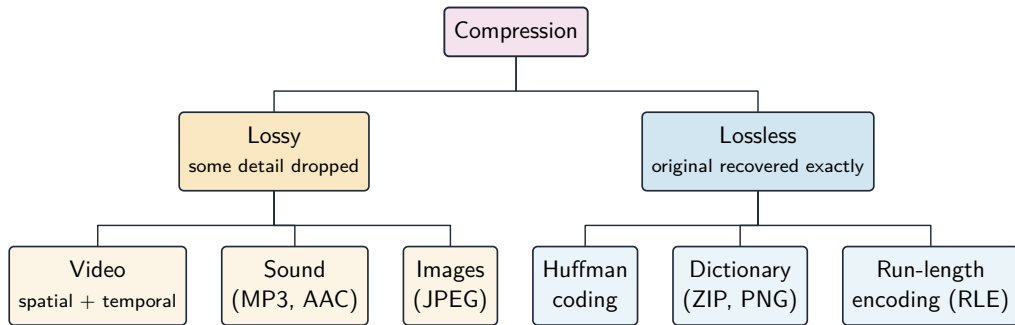
Binary Numbers



```
01000100 01000101 01010011 01010000 01001001 01000101 01010010 01010100 01000001 00100001 00100000 01001110
01001111 01010011 00100000 01000101 01010011 01010100 11000011 10000001 01001110 00100000 01001101 01000001
01001110 01001001 01010000 01010101 01001100 01000001 01001110 01000100 01001111 00100000 01000001 00100000
01010011 01010101 00100000 01000001 01001110 01010100 01001111 01001010 01001111 00101110 00100000 01000101
01001100 00100000 01001110 01010111 01001111 00100000 01001001 01001101 01010000 01001100 01001001 01000011
01000001 00100000 01001100 01000001 00100000 01010000 01000101 01010010 01000100 01001001 01000100 01000001
00100000 01000100 01000101 00100000 01010100 01001111 01000100 01000001 01010011 00100000 01001110 01010101
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01010100 01000001 01000100 01000101 01010011 00100000 01011001 00100000 01000101 01001100 00100000 01000011
01001111 01001110 01010100 01010010 01001111 01001100 00100000 01001101 01001001 01001110 01010101 01000011
01001001 01001111 01010011 01001111 00100000 01000100 01000101 00100000 01010100 01001111 01000100 01000001
00100000 01001100 01000001 00100000 01010000 01001111 01000010 01001100 01000001 01000011 01001001 11000011
10010011 01001110 00100000 01000001 01010011 11000011 10001101 00100000 01000011 01001111 01001101 01001111
00100000 01010101 01001110 01000001 00100000 01010010 01000101 01000100 01010101 01000011 01000011 01001001
11000011 10010011 01001110 00100000 01000100 01010010 11000011 10000001 01010011 01010100 01001001 01000011
01000001 00100000 01000100 01000101 00100000 01001100 01000001 00100000 01010000 01001111 01000010 01001100
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01000001 01001100 00100001 00100000 01010010 01000101 01010011 01001001 01010011 01010100 01000101 01001110
01000011 01001001 01000001 00100000 01001001 01001110 01010100 01000101 01001100 01000101 01000011 01010100
01010101 01000001 01001100 00100001 01100100 01100101 01110011 01110000 01100101 01110010 01110100 01100001
01100100 00100001 01100100 01100101 01110011 01110000 01100101 01110010 01110100 01100001 01100100 00100001
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01110011 01110000 01100101 01110010 01110100 01100001 01100100 00100001 01100100 01100101 01110011 01110000
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01110100 01100001 01100100 00100001 01100100 01100101 01110011 01110000 01100101 01110010 01110100 01100001
01100100 00100001 01100100 01100101 01110011 01110000 01100101 01110010 01100001 01100100 00100001
```

Binary digits on a display – all digital data is ultimately stored as 0s and 1s

Check yourself, in detail



Compression methods: lossless versus lossy, with common examples

Exam tips

- Convert confidently between **binary**, **decimal**, and (where asked) hexadecimal – practise until it is quick.
- Remember a bit is one binary digit and a byte is 8 bits; n bits represent 2^n values.
- Explain that all data – numbers, text, images, sound – is stored as binary, and that finite bits cause **overflow** and **round-off**.
- Distinguish **lossless** from **lossy** compression and when each is appropriate.
- Show the analog-to-digital idea: **sampling** turns a continuous signal into discrete values.

5

Recap

Unit 2 – key takeaways

1

Bases

binary, decimal, hex – until it is quick

2

Bits

n bits $\Rightarrow 2^n$ values; byte = 8 bits

3

Limits

overflow and round-off from finite bits

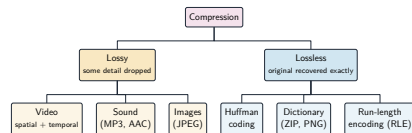
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Compression

lossless vs lossy; sampling analog to digital

Check yourself

- 1 Convert 1011_2 to decimal.
- 2 How many values can n bits represent?
- 3 Lossless or lossy for a program file – and why?
- 4 Does a correlation in the data prove causation?



Answers 1. 11 2. 2^n (from 0 to $2^n - 1$) 3. lossless – every bit must be restored exactly 4. no